

Appendix for ‘On Mates’ Puzzle’

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In this Appendix, I define two fragments of English and characterize their corresponding formal semantic computation. The first fragment captures individual *de re* belief reports. The second fragment covers Mates’ sentences in their opaque and transparent interpretation. The following aspects of the syntax and semantics are common to both fragments:

Syntax: Each expression is assigned a type among the following:

- u, e, s & t are types
- If a, b are types, $\langle a, b \rangle$ is a type

The language contains infinite variables v_a of each type a . I use the following conventions:

- x, y, z are of type e
- w, v, u are of type s
- G, G' are of type $\langle e, \langle s, e \rangle \rangle$
- φ, φ' are of type $\langle \langle e, \langle s, e \rangle \rangle, \langle s, t \rangle \rangle$
- $\mathcal{H}, \mathcal{H}'$ are of type $\langle u, \langle e, \langle s, t \rangle \rangle \rangle$
- ω, ω' are of type $\langle \langle u, \langle e, \langle s, t \rangle \rangle \rangle, \langle s, t \rangle \rangle$

Semantics: A semantic model is a tuple $M = \langle W, D_e, D_t, D_u, I \rangle$ where:

- W = a set of possible worlds
- D_e = a set of individuals
- D_t = a set of truth values $\{0, 1\}$
- D_u = a set of well-formed expressions, pairs of phonological and syntactic representations
- I = a valuation function mapping constant expressions to (sets of) elements of the various domains

In both fragments, the denotation of a sentence depends on:

- model M
- a context of utterance c
- a variable assignment g
- a possible world w

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Non-logical constants depend on the valuation function I provided by the model; variables depend on the assignment function g . To combine the meanings of expressions, we use Functional Application (FA):

- If α is of type $\langle a, b \rangle$, β is of type a and γ is an expression consisting of α and β , then: $\gamma = \alpha(\beta)$

Functional expressions are represented using lambda-abstraction.

De re about individuals

In the fragment for *de re* reports about individuals, the denotation of a sentence depends on the above parameters and:

- a constant function $\Pi = \lambda x \lambda w . x$

In this fragment, sentences are uniformly functions from objects of type $\langle e, \langle s, e \rangle \rangle$ to propositions (type $\langle s, t \rangle$). To interpret any sentence at an index of evaluation, the index needs to provide more than a possible world. That is the role of Π : it provides an argument for sentences like ‘Ortcutt is a spy’ whereby, if this clause contains a CG-variable as sister to ‘Ortcutt’, Π maps ‘Ortcutt’ to its transparent value, namely Ortcutt. Otherwise Π is idle.

Where α is a sentence, we write $\llbracket \alpha \rrbracket^{M,c,g,\Pi,w}$ for ‘the denotation of α with respect to M, c, g, Π, w ’. However, since M, c, g do not play a role in our derivation, we simply write $\llbracket \alpha \rrbracket^{\Pi,w}$.

Constants have the following denotation (assigned by function I):

- Ralph = r
- Ortcutt = o
- that = $\lambda \varphi . \varphi$
- is-a-spy = $\lambda x \lambda G \lambda w . \mathbf{spy}(x)(w)$ (*de dicto*)
- is-a-spy = $\lambda x \lambda G \lambda w . \mathbf{spy}(G(x)(w))(w)$ (*de re*)
- believes = $\lambda \varphi \lambda x \lambda G \lambda w . \exists G'(G' \text{ is an } ACG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \varphi(G')(v))$

The predicate ‘is-a-spy’ is ambiguous between its *de re* and *de dicto* interpretation. In its *de re* interpretation, it contains a CG-variable that scopes over its first argument.

The sentence (2a) ‘Ralph believes that Ortcutt is a spy’ has the following LF (in its *de re* interpretation),

(2a) Ralph [believes- w_0 [that [λG_1 [λw_1 [[[[[G_1 Ortcutt] w_1] is a spy] w_1]]]]]]

which has the following compositional derivation (all transformations are the result of FA):

- believes(that(is-a-spy(Ortcutt)))(Ralph) =
- $\lambda \varphi \lambda x \lambda G \lambda w . \exists G'(G' \text{ is an } ACG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \varphi(G')(v))$
 $(\lambda \varphi . \varphi(\lambda x \lambda G \lambda w . \mathbf{spy}(G(x)(w))(w)(o)))(r) =$
- $\lambda \varphi \lambda x \lambda G \lambda w . \exists G'(G' \text{ is an } ACG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \varphi(G')(v))$
 $(\lambda \varphi . \varphi(\lambda G \lambda w . \mathbf{spy}(G(o)(w))(w)))(r) =$
- $\lambda \varphi \lambda x \lambda G \lambda w . \exists G'(G' \text{ is an } ACG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \varphi(G')(v))(\lambda G \lambda w . \mathbf{spy}(G(o)(w))(w))(r) =$

- $\lambda x \lambda G \lambda w. \exists G' (G' \text{ is an } ACG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \lambda G \lambda w. \mathbf{spy}(G(o)(w))(w)(G')(v))(r) =$
- $\lambda G \lambda w. \exists G' (G' \text{ is an } ACG_{r,w} \ \& \ \forall v \in DOX_{r,w} : \lambda G \lambda w. \mathbf{spy}(G(o)(w))(w)(G')(v)) =$
- $\lambda G \lambda w. \exists G' (G' \text{ is an } ACG_{r,w} \ \& \ \forall v \in DOX_{r,w} : \mathbf{spy}(G'(o)(v))(v))$

This sentence is interpreted as follows:

- $\llbracket (2a) \rrbracket^{\Pi,w} = 1$ iff $\exists G' (G' \text{ is an } ACG_{r,w} \ \& \ \forall v \in DOX_{r,w} : \mathbf{spy}(G'(o)(v))(v))$.

In words: there exists an acquaintance-based concept generator G' for Ralph at w such that, across Ralph's doxastic alternatives, G' maps Orcutt to a spy.

De re about words

In this fragment, 'believe' maps *metalinguistic* concept-generators to propositions. Correspondingly, all other expressions have a different semantic type as well. The denotation of a sentence depends on M, c, g, w and:

- a constant disquotational function Δ (see **Definition 6** and the surrounding discussion)

Where α is a sentence, we write $\llbracket \alpha \rrbracket^{M,c,g,\Delta,w}$, or $\llbracket \alpha \rrbracket^{\Delta,w}$ for short.

Constants have the following denotation (assigned by I):

- Andrea = a
- Fiona = f
- that = $\lambda \omega. \omega$
- not = $\lambda \omega \lambda \mathcal{H} \lambda w. \mathbf{not} \ \omega(\mathcal{H})(w)$
- is-an-attorney = $\lambda x \lambda \mathcal{H} \lambda w. \mathcal{H}(\mathit{attorney})(x)(w)$ *(opaque)*
- is-a-lawyer = $\lambda x \lambda \mathcal{H} \lambda w. \mathcal{H}(\mathit{lawyer})(x)(w)$ *(opaque)*
- is-an-attorney = is-a-lawyer = $\lambda x \lambda \mathcal{H} \lambda w. \mathbf{lawyer}(x)(w)$ *(transparent)*
- believes = $\lambda \omega \lambda x \lambda \mathcal{H} \lambda w. \exists \mathcal{H}' (\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$

The predicates 'is-an-attorney' and 'is-a-lawyer' are ambiguous between *transparent* and *opaque* interpretations. In the former case, they have the same denotation; in the latter, they do not: each contains an MCG-variable as sister to a different member of D_u , possibly resulting in different denotations.

The sentence (1a) 'Andrea believes that Fiona is an attorney' has the following LF (in its *opaque* interpretation):

(1a) Andrea [believes- w_0 [that [$\lambda \mathcal{H}_1$ [λw_1 [w_1 [Fiona [\mathcal{H}_1 is-an-attorney]]]]]]]

And the following compositional derivation:

- believes(that(is-an-attorney(Fiona)))(Andrea) =
- $\lambda \omega \lambda x \lambda \mathcal{H} \lambda w. \exists \mathcal{H}' (\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda \omega. \omega(\lambda x \lambda \mathcal{H} \lambda w. \mathcal{H}(\mathit{attorney})(x)(w)(f)))(a) =$

- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda\mathcal{H}\lambda w.\mathcal{H}(\text{attorney})(f)(w)))(a) =$
- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\mathcal{H}\lambda w.\mathcal{H}(\text{attorney})(f)(w))(a) =$
- $\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \lambda\mathcal{H}\lambda w.\mathcal{H}(\text{attorney})(f)(w)(\mathcal{H}')(v))(a) =$
- $\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \lambda\mathcal{H}\lambda w.\mathcal{H}(\text{attorney})(f)(w)(\mathcal{H}')(v)) =$
- $\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \mathcal{H}'(\text{attorney})(f)(v)).$

This sentence is interpreted as follows:

- $\llbracket(1a)\rrbracket^{\Delta,w} = 1$ iff $\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \mathcal{H}'(\text{attorney})(f)(v))$

In words: there exists a metalinguistic concept generator for Andrea at w that maps ‘attorney’ ($\in D_u$) to a property that Fiona has across Andrea’s doxastic alternatives. In Context A, the truth of this sentence is witnessed by a metalinguistic concept generator that maps ‘attorney’ to the property of *being a judge*.

The sentence (1b) (‘Andrea believes that Fiona is not a lawyer’) has the following LF (in its *opaque* interpretation) and derivation:

(1b) Andrea [believes- w_0 [that [not [$\lambda\mathcal{H}_1$ [λw_1 [w_1 [Fiona [\mathcal{H}_1 is-a-lawyer]]]]]]]]

- believes(that(not(is-a-lawyer(Fiona)))(Andrea)=
- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda\omega\lambda\mathcal{H}\lambda w.\mathbf{not} \ \omega(\mathcal{H})(w)(\lambda x\lambda\mathcal{H}\lambda w.\mathcal{H}(\text{lawyer})(x)(w)(f))))(a) =$
- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda\omega\lambda\mathcal{H}\lambda w.\mathbf{not} \ \omega(\mathcal{H})(w)(\lambda\mathcal{H}\lambda w.\mathcal{H}(\text{lawyer})(f)(w))))(a) =$
- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda\mathcal{H}\lambda w.\mathbf{not} \ \lambda\mathcal{H}\lambda w.\mathcal{H}(\text{lawyer})(f)(w)))(\mathcal{H})(w))(a) =$
- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda\mathcal{H}\lambda w.\mathbf{not} \ \mathcal{H}(\text{lawyer})(f)(w)))(a) =$
- $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\mathcal{H}\lambda w.\mathbf{not} \ \mathcal{H}(\text{lawyer})(f)(w))(a) =$
- $\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \lambda\mathcal{H}\lambda w.\mathbf{not} \ \mathcal{H}(\text{lawyer})(f)(w)(\mathcal{H}')(v))$
 $(a) =$
- $\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \lambda\mathcal{H}\lambda w.\mathbf{not} \ \mathcal{H}(\text{lawyer})(f)(w)(\mathcal{H}')(v)) =$
- $\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \mathbf{not} \ \mathcal{H}'(\text{lawyer})(f)(v))$

This sentence is interpreted as follows:

- $\llbracket(1b)\rrbracket^{\Delta,w} = 1$ iff $\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \mathbf{not} \ \mathcal{H}'(\text{lawyer})(f)(v))$

In words: there exists a metalinguistic concept generator for Andrea at w that maps ‘lawyer’ ($\in D_u$) to a property that Fiona lacks across Andrea’s doxastic alternatives. In Context A, the truth of this sentence is witnessed by a metalinguistic concept generator that maps ‘lawyer’ to the property of *being a lawyer*, which Andrea knows that Fiona does not have.

Finally, in its *transparent* interpretation, (1a) has the following LF and derivation:

- (1a) Andrea [believes- w [that [$\lambda\mathcal{H}_1$ [λw_1 [w_1 [Fiona is an attorney]]]]]
- believes(that(is-an-attorney(Fiona)))(Andrea)=
 - $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda x\lambda\mathcal{H}\lambda w.\mathbf{lawyer}(x)(w)(f)))(a) =$
 - $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\omega.\omega(\lambda\mathcal{H}\lambda w.\mathbf{lawyer}(f)(w)))(a) =$
 - $\lambda\omega\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \omega(\mathcal{H}')(v))$
 $(\lambda\mathcal{H}\lambda w.\mathbf{lawyer}(f)(w))(a) =$
 - $\lambda x\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{x,w} \ \& \ \forall v \in DOX_{x,w} : \lambda\mathcal{H}\lambda w.\mathbf{lawyer}(f)(w)(\mathcal{H}')(v))(a) =$
 - $\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \lambda\mathcal{H}\lambda w.\mathbf{lawyer}(f)(w)(\mathcal{H}')(v)) =$
 - $\lambda\mathcal{H}\lambda w.\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \mathbf{lawyer}(f)(v))$

This sentence is interpreted as follows:

- $\llbracket(1a)\rrbracket^{\Delta,w} = 1$ iff $\exists\mathcal{H}'(\mathcal{H}' \text{ is an } MCG_{a,w} \ \& \ \forall v \in DOX_{a,w} : \mathbf{lawyer}(f)(v))$

In words: there exists an MCG for Andrea at w and Fiona is a lawyer across Andrea's doxastic alternatives.