

# Dynamic inquisitive semantics

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- ① To develop a logical framework that combines insights from:
  - Dynamic semantics
  - Inquisitive semantics
- ② To highlight some of the advantages that such a framework has w.r.t. frameworks that are non-dynamic or non-inquisitive in the analyses of questions.

- 1 Motivating dynamic inquisitive semantics
- 2 Dynamic inquisitive semantics
- 3 Explaining the motivating data

1 Motivating dynamic inquisitive semantics

2 Dynamic inquisitive semantics

3 Explaining the motivating data

# Motivation for a dynamic semantics of questions

- A dynamic semantics of questions is needed to capture certain types of **anaphora**.

(Groenendijk, 1998; van Rooij, 1998; Haida, 2007)

- It also provides an attractive account of certain kinds of **intervention effects**.

(Haida, 2007)

# Anaphora

- In dynamic semantics, meaning is not equated with truth-conditions but rather with **context change potential**.
- This includes the potential to introduce **discourse referents**.
- Dynamic semantics has been motivated by contrasts like:
  - (1) a. [One of my ten marbles]<sup>x</sup> is not here.  
It<sub>x</sub> is probably under the sofa.
  - b. Nine of my ten marbles are here.  
#It is probably under the sofa.
- The first sentence in (1-a) introduces a discourse referent picked up by the anaphoric pronoun.
- The first sentence in (1-b) is truth-conditionally equivalent but does not introduce a suitable discourse referent.

## Anaphora in questions

- Dynamic semantics for questions is motivated by examples like:
  - (2) a. Which<sup>x</sup> one of your three sons is living in Paris?  
And is he<sub>x</sub> happy there?
  - b. Which<sup>x</sup> two of your three sons are not living in Paris? #And is he<sub>x</sub> happy there?
- The initial questions in (2-a) and (2-b) are equivalent in terms of **resolution/answerhood conditions**.
- But they differ in their potential to license pronominal anaphora.

## Anaphora in questions

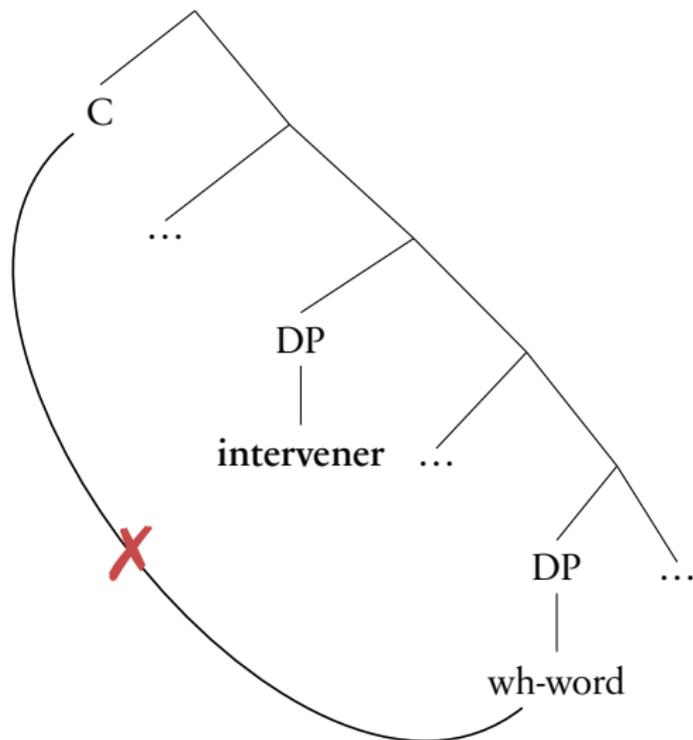
- The fact that wh-questions license anaphora has been the primary piece of motivation for dynamic approaches to question semantics.

(Groenendijk, 1998; van Rooij, 1998; Haida, 2007)

- Wh-words are taken to introduce discourse referents, just like plain existential indefinites, which can be picked up by anaphoric pronouns.

# Intervention effects

Certain operators appearing between a wh-word and the interrogative complementizer lead to ungrammaticality.



## Intervention effects: an example

Example from Beck (2006):

- (3) a. Wer hat **Luisse** **wo** angetroffen?  
who-Nom has Luise where met  
'Who met Luise where?'
- b. ??Wer hat **niemandem** **wo** angetroffen?  
who-Nom has nobody-Dat where met  
'Who didn't meet anybody where?'
- c. Wer hat **wo** **niemandem** angetroffen?  
who-Nom has where nobody-Dat met  
'Who didn't meet anybody where?'

## Intervention effects: possible interveners

- Possible interveners (Beck, 2006):
  - Focus sensitive operators: *only, even,...*
  - Nominal quantifiers: *every, no, most, few,...*
  - Adverbial quantifiers: *always, often, never,...*
  - Negation: *not*

## Intervention effects: two approaches

- How to understand the licensing relation between the wh-word and the complementizer?
- Two prominent approaches (among others):
  - **Focus approach** (Beck, 2006) – wh-words introduce focus alternatives. Intervention effects arise when C cannot access these focus alternatives because they are consumed by a focus sensitive intervener.
  - **Dynamic approach** (Haida, 2007) – wh-words introduce discourse referents that C has to access. Intervention effects arise if such access is blocked by operators that do not let discourse referents project from their scope.

## Intervention effects: the focus approach

- The **focus approach** works well for **only** and **even**, which are without doubt focus sensitive.
- However, nominal quantifiers **every**, **no**, and **most** are more problematic because they are not necessarily focus sensitive. Beck (2006, Section 4), and Haida (2007, Chapter 8)
- Experimental data suggest that the focus-sensitive particle **also** is not an intervener in German. Haida and Repp (2013)

## Intervention effects: the dynamic approach

- The **dynamic approach** works well for **quantifiers** and **negation**, which are known to block discourse referents from projecting:
  - (4) a. John didn't consider buying [a car]<sup>i</sup>.  
\*It<sub>i</sub> was too expensive.
  - b. Most students considered buying [a car]<sup>i</sup>.  
\*It<sub>i</sub> was very cheap.
  - c. John often considered buying [a car]<sup>i</sup>.  
\*It<sub>i</sub> was very cheap.
- **Focus sensitive particles** do not block discourse referents from projecting, so require a different explanation (Haida, 2007)
- **Cross-linguistic variation** as to which operators act as interveners is largely an open issue for both approaches.

# Motivation for an inquisitive semantics of questions

- Dynamic theories of questions are based on **partition semantics** (Groenendijk, 1998; van Rooij, 1998; Haida, 2007)
- Partition semantics is suitable to model **exhaustive** interpretations of questions:
  - (5) Which of the guests are vegetarian?

## Motivation for an inquisitive semantics of questions

- However, **non-exhaustive** question interpretations are difficult to capture in partition semantics

(6) Who has a bike that I could borrow for 15 minutes?

- In **inquisitive semantics** both **exhaustive** and **non-exhaustive** question interpretations can be captured straightforwardly
- Moreover, **dynamic inquisitive semantics** allows us to relate exhaustive and non-exhaustive questions to another contrast, strong and weak readings of donkey anaphora

## Summary

	Anaphora	Intervention	Non-exhaustive
Static inquisitive	No	No	Yes
Dynamic partition	Yes	Yes	No
<b>Dynamic inquisitive</b>	<b>Yes</b>	<b>Yes</b>	<b>Yes</b>

- There are other potential advantages to dynamic inquisitive semantics which we do not explicitly discuss here

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# Dynamic inquisitive semantics

Our compositional dynamic inquisitive semantics ( $\text{Inq}_D$ ) integrates the basic static inquisitive system,  $\text{Inq}_B$ , with

- the dynamic system GSV (Groenendijk *et al.*, 1996), and
- the compositional dynamic system CDRT (Muskens, 1996)

The system presented here simplified.

$\text{Inq}_D$  for a first-order language:

to appear in Dotlačil and Roelofsen, Proceedings of SuB.

Contexts in GSV encapsulate:

- 1 information about the world
- 2 information about the discourse referents
- 3 information about dependencies between the world and the discourse referents

Formally, a context in GSV is a set  $s$  of  $\langle w, g \rangle$  pairs (possibilities), such that  $w$  is a possible world and  $g$  a (partial) assignment function

This encodes:

- Information about the world:

$$\mathbf{worlds}(s) := \{w \mid \langle w, g \rangle \in s \text{ for some } g\}$$

- Information about the discourse referents:

$$\mathbf{assigns}(s) := \{g \mid \langle w, g \rangle \in s \text{ for some } w\}$$

- Information about dependencies between the world and the discourse referents

Limits on the notion of context in GSV:

- Contexts represent information, but not **contextual issues**
- Contexts in GSV are not rich enough to analyze questions

# Contexts in inquisitive semantics

Contexts in inquisitive semantics encapsulate:

- 1 information about the world
- 2 issues raised about the world

# Contexts in inquisitive semantics

Context  $C$  – a set of information states, each a set of possible worlds

Conditions on  $C$ :

- each information state contains enough information to resolve the raised issues
- no information state contains any worlds that have been ruled out by the available information
- set of information states is downward closed and non-empty
- the information available in a context,  $\text{INFO}(c) := \bigcup c$ .
- A context  $c$  is inquisitive just in case  $\text{INFO}(c) \notin c$ .

# Contexts in inquisitive semantics

Limits on the notion of context in  $\text{Inq}_B$ :

- Contexts represent information and issues about the world
- They do not model information and issues about discourse referents
- They do not model dependencies between the world and the discourse referents

# Contexts in dynamic inquisitive semantics

Context  $C$  – a set of information states, each a set of possibilities  
( $\langle w, g \rangle$  pairs)  
( $w$  – world;  $g$  – dref assignment function)

Conditions on  $C$ :

- each information state contains enough information to resolve the raised issues
- no information state contains any possibilities that have been ruled out by contextual information
- set of information states is downward closed and non-empty

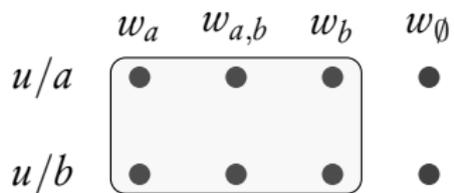
## Contexts in dynamic inquisitive semantics – conventions

- $\text{Inq}_D$  has four basic types:  $e$  for individuals,  $s$  for possible worlds,  $t$  for truth values, and  $r$  for discourse referents

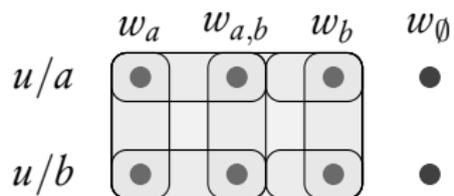
Object	Type	Type abbreviation	Variable convention
dref assignment f.	$(re)$	$a$	-
possibility	$(s \times a)$	-	$p$
information state	$((s \times a)t)$	$i$	$s$
context	$(it)$	$k$	$c, c'$
update function	$(kk)$	$T$	$A, B$

**Table:** Types and abbreviation conventions

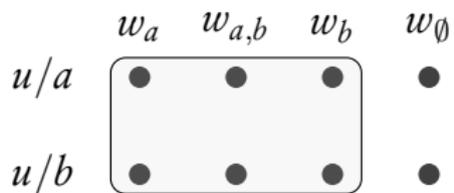
# Contexts in dynamic inquisitive semantics – graphical representation



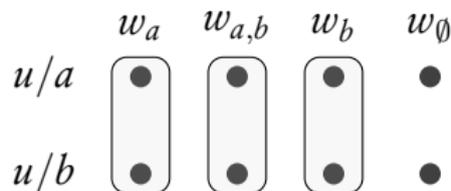
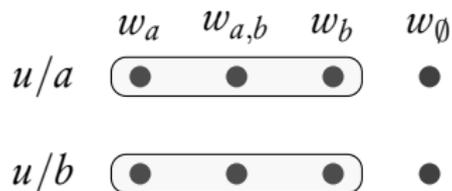
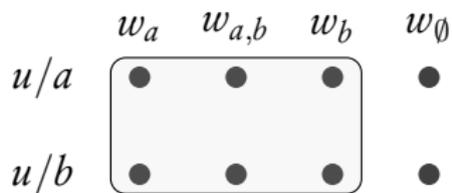
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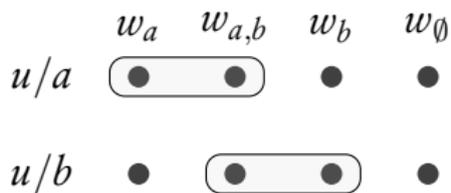
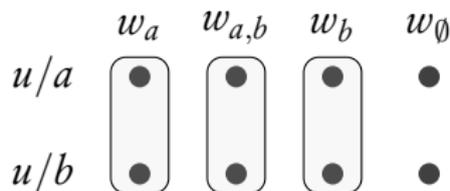
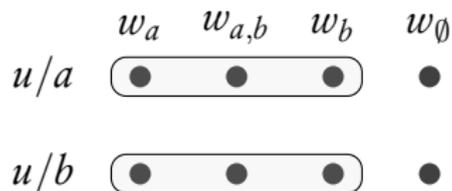
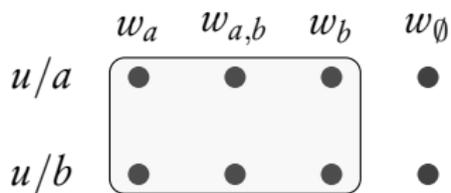
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# Contexts in dynamic inquisitive semantics – graphical representation



## Extension and subsistence

A state  $s'$  **extends** a state  $s$ ,  $s' \geq s$ , iff  $s' = s$  or:

- $s'$  extends the information about the **world** provided in  $s$  (by excluding some possibilities), and/or
- $s'$  extends the information about the **discourse referents** provided in  $s$  (by adding new discourse referents)

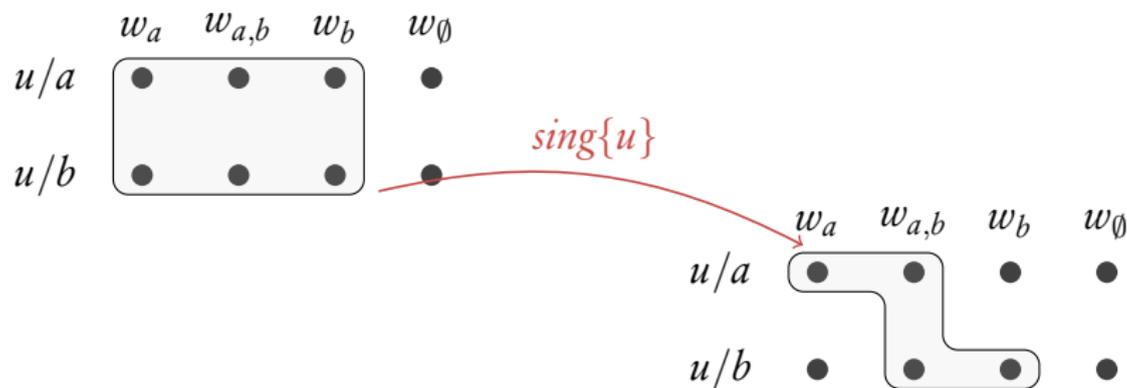
A state  $s$  **subsists** in a state  $s'$  iff:

- $s' \geq s$ , and
- every possibility in  $s$  is in  $s'$ , modulo the addition of new discourse referents

A state  $s$  subsists in a context  $C$  iff there is some  $s' \in C$  such that  $s$  subsists in  $s'$ . In this case,  $s'$  is called a descendant of  $s$ .

# Semantics: relations and conjunction

$$(7) \quad R\{u_1, \dots, u_n\} := \lambda c_k \lambda s_i. \quad s \in c \wedge \forall p \in s \\ (R(w_p)(g_p(u_1)) \dots (g_p(u_n)))$$



$$(8) \quad A_{(kk)}; B_{(kk)} := \lambda c_k. B(A(c))$$

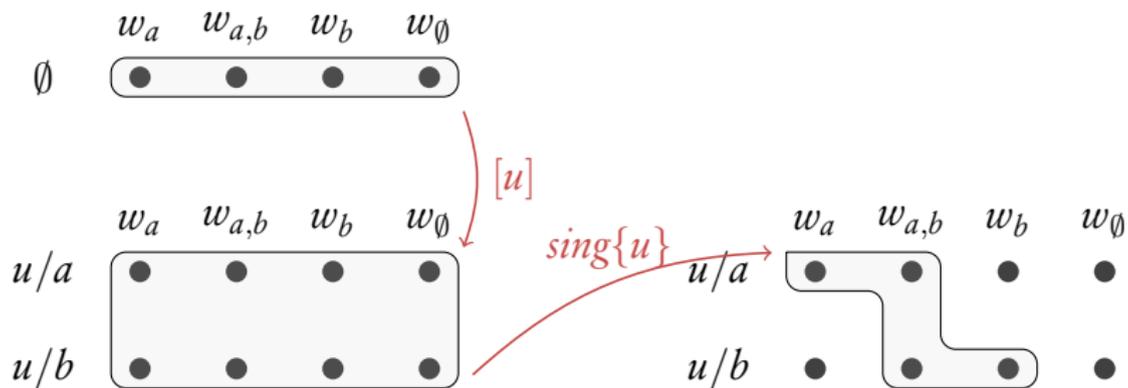
## Semantics: dref introduction

- If  $C$  is a context and  $u$  not in assignments in  $C$ , then introducing the dref  $u$  creates the largest context  $C'$  such that every  $s' \in C'$  has some  $s \in C$  s.t.  $s$  subsists in  $s'$  and  $\mathbf{assigns}(s') = \mathbf{assigns}(s) \cup \{u\}$ .

## Semantics: dref introduction

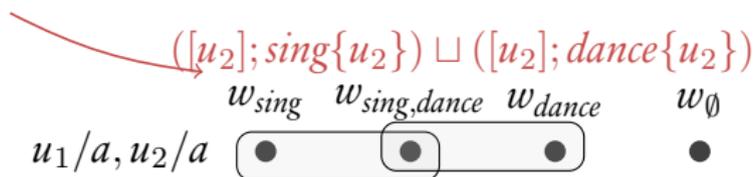
- If  $C$  is a context and  $u$  not in assignments in  $C$ , then introducing the dref  $u$  creates the largest context  $C'$  such that every  $s' \in C'$  has some  $s \in C$  s.t.  $s$  subsists in  $s'$  and  $\mathbf{assigns}(s') = \mathbf{assigns}(s) \cup \{u\}$ .
- $\llbracket g[u]g' \rrbracket^{F,I,\theta} = 1$  iff
  - $\llbracket \mathbf{dom}(g') \rrbracket^{F,I,\theta} = \llbracket \mathbf{dom}(g) \rrbracket^{F,I,\theta} \cup \{\llbracket u \rrbracket^{F,I,\theta}\}$  and
  - $\llbracket \forall v_r((v \neq u \wedge v \in \mathbf{dom}(g)) \rightarrow g(v) = g'(v)) \rrbracket^{F,I,\theta} = 1$
- $\llbracket p[u]p' \rrbracket^{F,I,\theta} = 1$  iff  $\llbracket w_p = w_{p'} \rrbracket^{F,I,\theta} = 1$  and  $\llbracket g_p[u]g_{p'} \rrbracket^{F,I,\theta} = 1$
- $[u] := \lambda c_k \lambda s_i. \left\{ \begin{array}{l} \exists s' \in c. \\ \forall p \in s. \exists p' \in s'. (p'[u]p) \wedge \\ \forall p' \in s'. \exists p \in s. (p'[u]p) \end{array} \right\}$

# Semantics: dref introduction



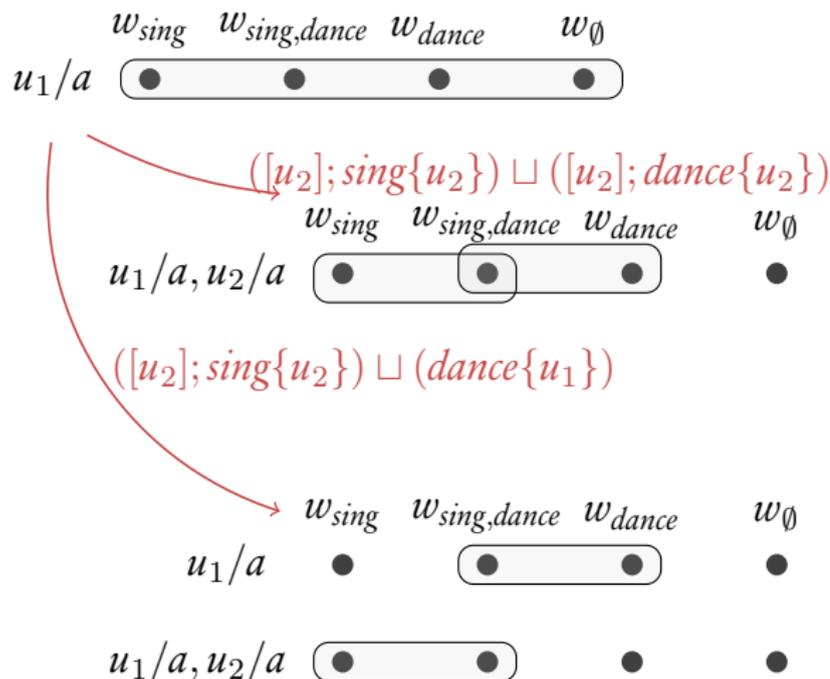
# Semantics: disjunction

$$(9) \quad A_{(kk)} \sqcup B_{(kk)} := \lambda c_k. A(c) \cup B(c)$$



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$$(9) \quad A_{(kk)} \sqcup B_{(kk)} := \lambda c_k. A(c) \cup B(c)$$



## Semantics: disjunction

- (10) Bill either rented a<sup>u</sup> car or hitchhiked.  
\*It<sub>u</sub> was probably a cabriolet.
- (11) Bill either rented a<sup>u</sup> blue car or a<sup>u</sup> red car.  
It<sub>u</sub> was probably a cabriolet.

cf. Stone (1992)

## Semantics: disjunction

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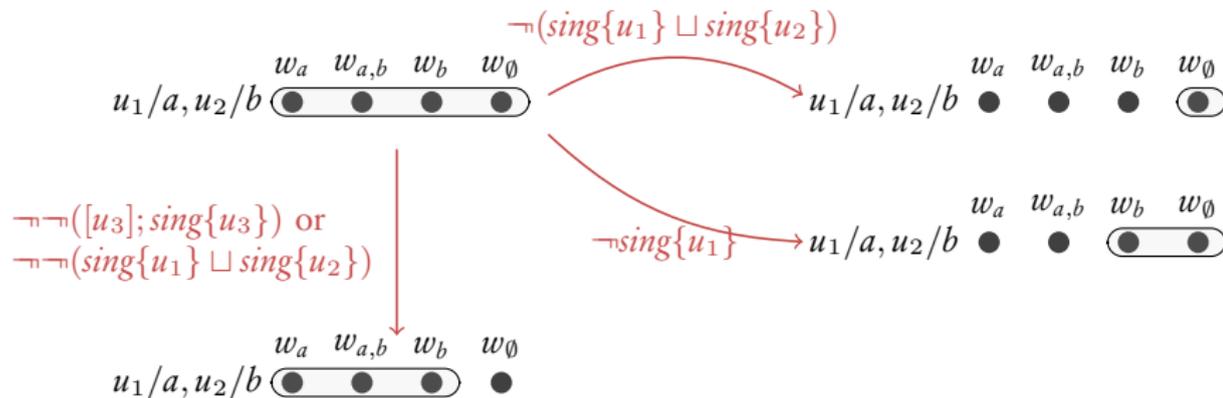
(11) Bill either rented a<sup>u</sup> blue car or a<sup>u</sup> red car.  
It<sub>u</sub> was probably a cabriolet.

cf. Stone (1992)

- (12) **A:** Bill either rented a<sup>u</sup> car or hitchhiked.  
**B:** The former, of course. It<sub>u</sub> was a cabriolet.
- (13) **A:** Did Bill rent a<sup>u</sup> car<sup>↑</sup> or did he hitchhike?<sup>↓</sup>  
**B:** The former, of course. It<sub>u</sub> was a cabriolet.

# Semantics: negation

$$(14) \quad \neg A_{(kk)} := \lambda c_k \lambda s_i. s \in c \wedge \neg \exists t \subseteq s (t \neq \emptyset \wedge t \text{ subsists in } A(c))$$



## Semantics: discharging inquisitiveness

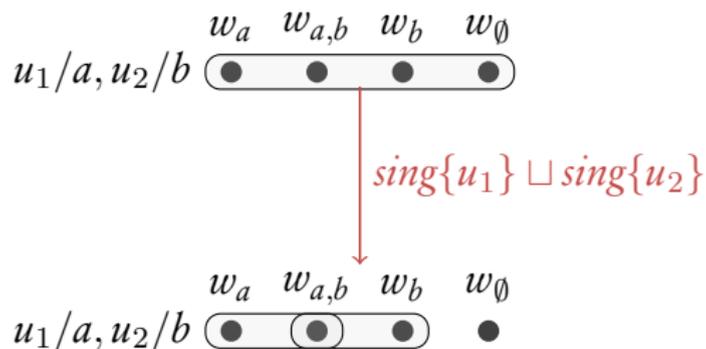
(15)  $!A_{(kk)} := \lambda c_k \lambda s_I. s \in !(A(c)) \wedge \exists s' \in c (s \text{ is an extension of } s')$

- For any context  $C$ :

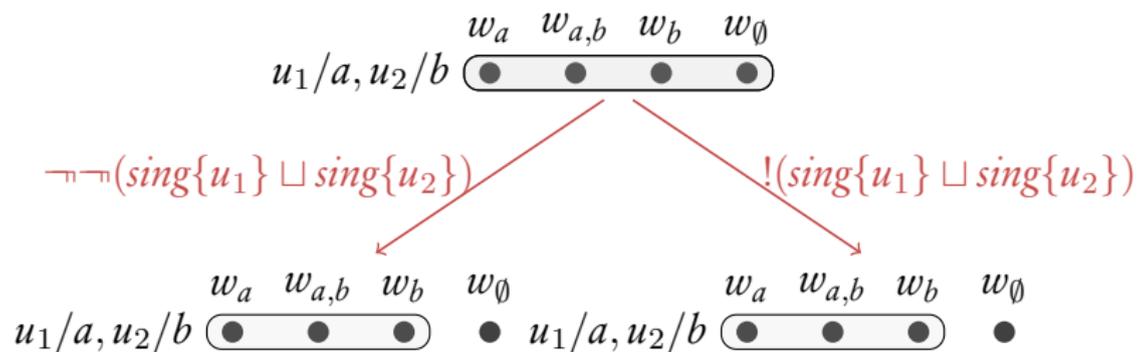
$!C := \{\bigcup C\}^\downarrow$  is the non-inquisitive projection of  $C$

- $!A$  discharges any inquisitiveness introduced by  $A$

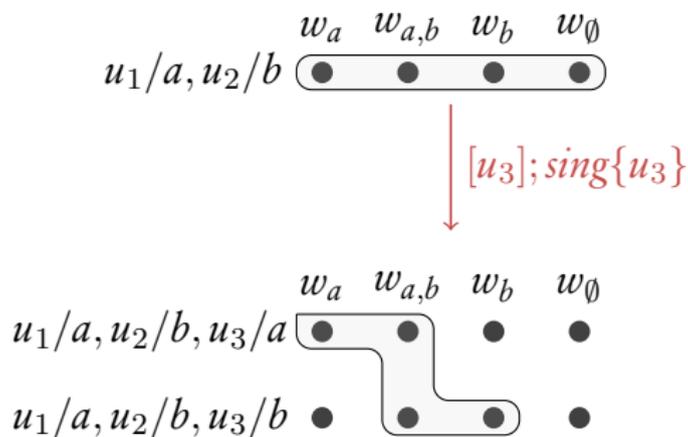
# Semantics: negation and exclamation operator (graphical comparison)



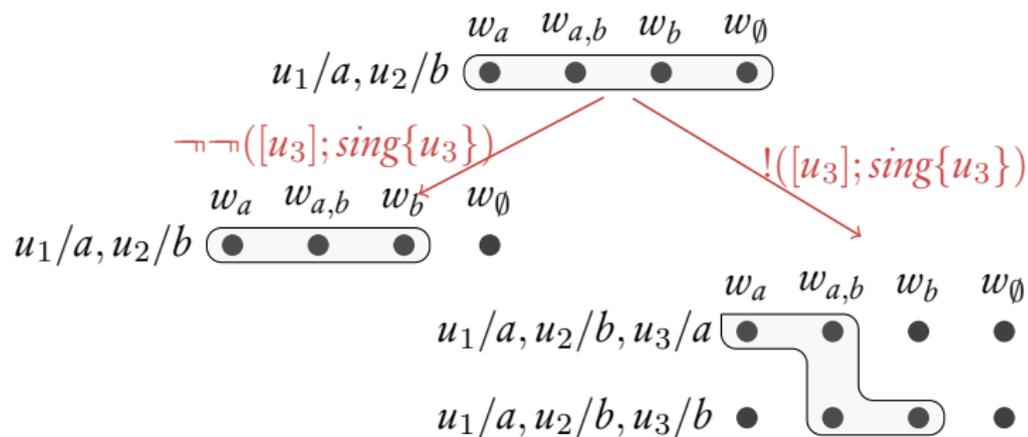
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# Semantics: ensuring inquisitiveness

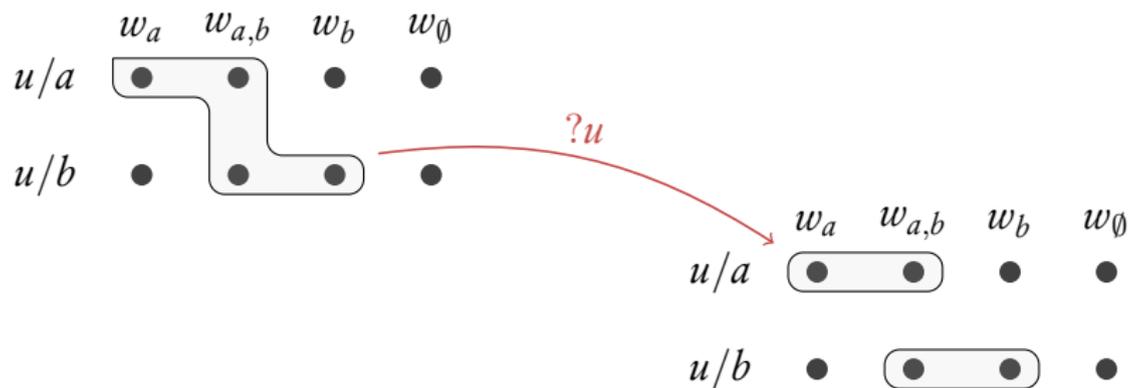
$$(16) \quad ?A_{(kk)} := A \sqcup \neg A$$

$$(17) \quad \langle ? \rangle A := \begin{cases} ?A & \text{if } A \text{ is not inquisitive} \\ A & \text{otherwise} \end{cases}$$



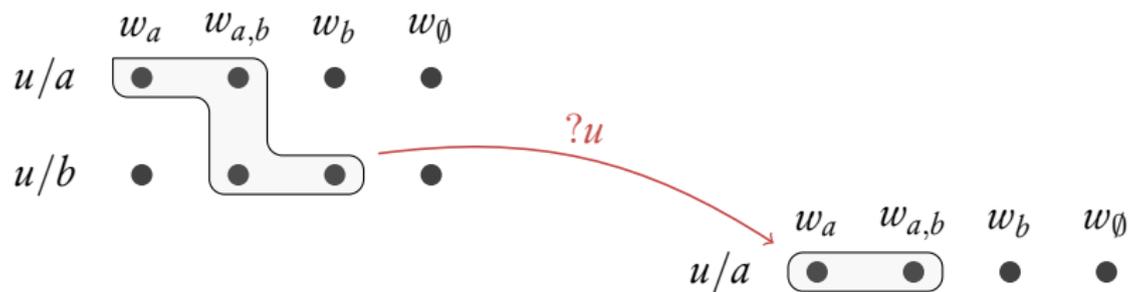
# Semantics: identity of drefs

$$(18) \quad ?u := \lambda c_k \lambda s_i. s \in c \wedge \exists x_e. \forall p \in s. g_p(u) = x$$



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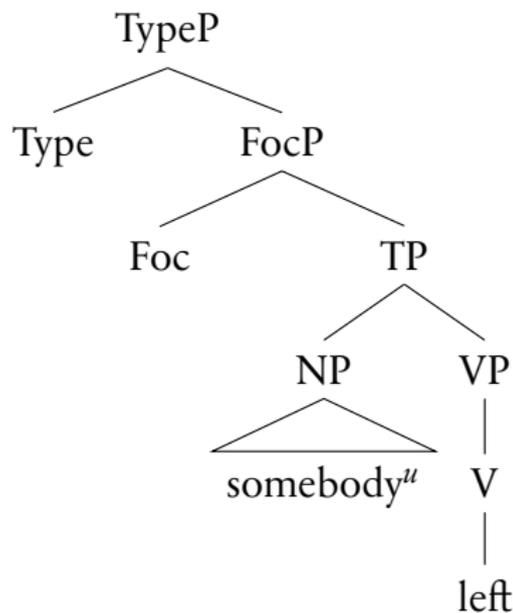


$$(19) \quad ?u_1, \dots, u_j, u_n := \lambda c_k \lambda s_i.$$



$$\left\{ \begin{array}{l} s \in c \wedge \\ \exists f. \forall p \in s. g_p(u_n) = f(g_p(u_1), \dots, g_p(u_j)) \end{array} \right\}$$

# Translation of declaratives



$$\llbracket \text{Foc} \rrbracket = \lambda A_{(kk)}.!(A)$$

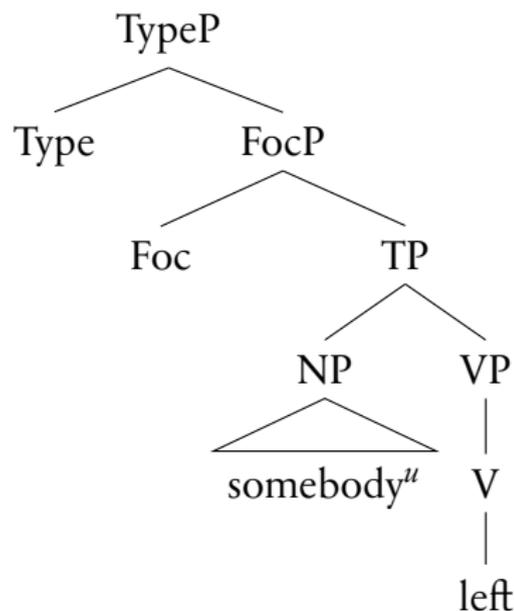
removes inquisitiveness of complement

$$\llbracket \text{leave} \rrbracket = \lambda v_r. \text{leave}\{v\}$$

$$\llbracket \text{somebody}^u \rrbracket = \lambda P_{rT}. [u]; P(u)$$

$$(20) \quad \llbracket \llbracket \text{FocP} \text{Somebody}^u \text{ left} \rrbracket \rrbracket = !([u]; \text{left}\{u\})$$

# Translation of declaratives

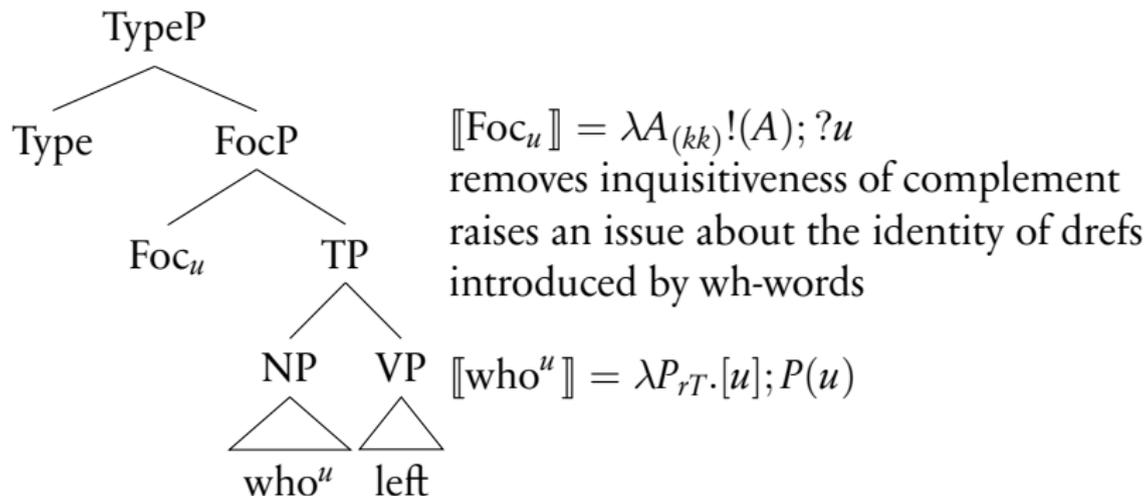


$\llbracket \text{Type} \rrbracket = \lambda A_{(kk)}.!(A)$   
removes inquisitiveness of complement

(21)  $\llbracket \text{Somebody}'' \text{ left} \rrbracket = !([u]; \text{left}\{u\})$

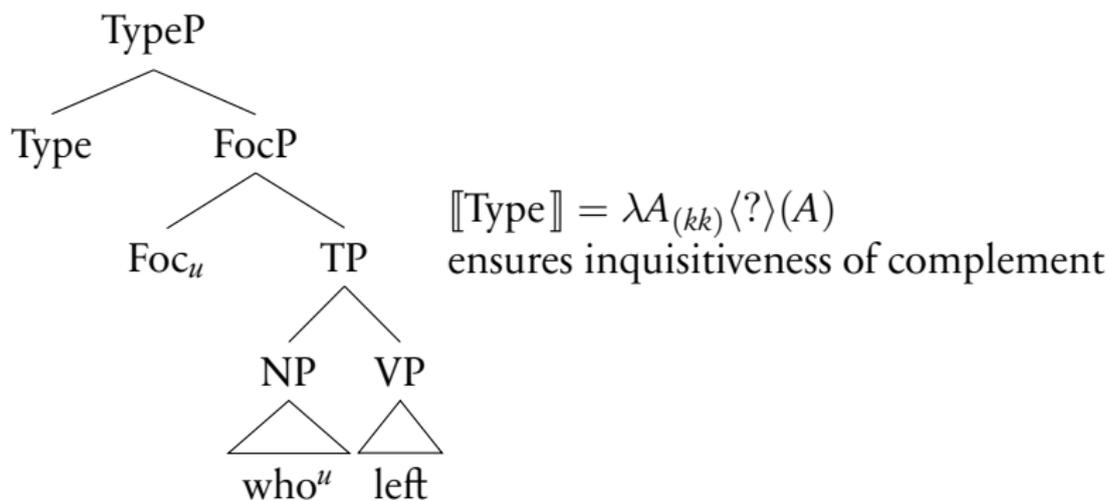
- (21) introduces a dref that left but (21) does not raise issues

## Translation of wh examples



$$(22) \quad \llbracket [\text{FocP Who}^u \text{ left}] \rrbracket = !([u]; \text{left}\{u\}); ?u$$

## Translation of wh examples

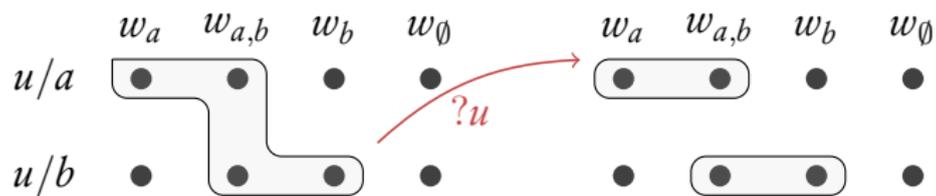


(23)  $\llbracket \text{Who}^u \text{ left} \rrbracket = !([u]; \text{left}\{u\}); ?u$

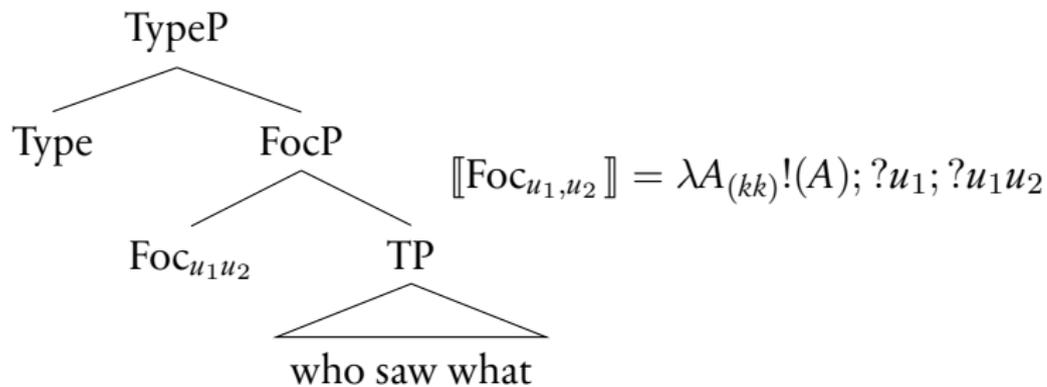
- (23) introduces a dref that left and (23) raises an issue about the identity of the dref

# Translations of wh examples

(24)  $\llbracket \text{Who left} \rrbracket = !([u]; \text{left}\{u\}); ?u$



## Translations of wh examples



(25)  $\llbracket \text{Who saw what} \rrbracket = !([u_1]; [u_2]; \text{saw}\{u_1, u_2\}); ?u_1; ?u_1 u_2$

## Explaining anaphora

(26) [[Someone<sup>u</sup> left. He<sub>u</sub> was furious.]]  
= !([u]; left{u}); !(furious{u})

(27) [[Who<sup>u</sup> left? (I don't know but) he<sub>u</sub> was furious.]]  
= !([u]; left{u}); ?u; !(furious{u})

- Wh-words are indefinites.
- Their inquisitiveness is introduced by the Foc head that they agree with.
- This does not affect their binding possibilities.
- Their licensing of anaphora is captured.

## Explaining anaphora

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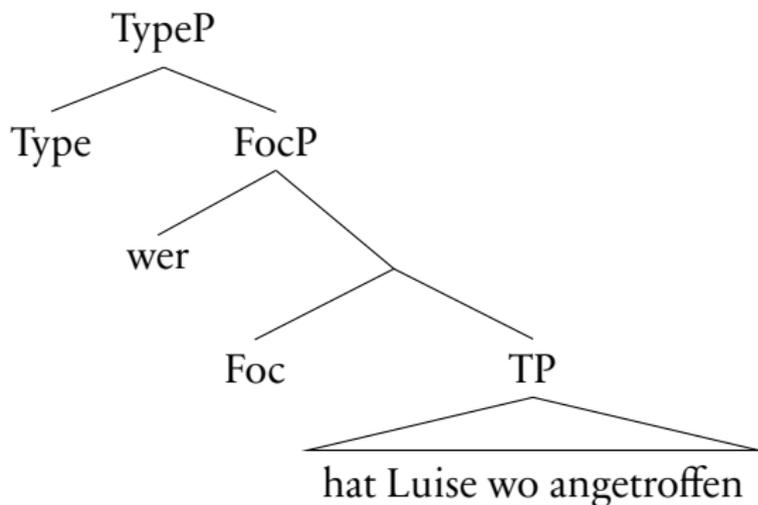
(27) [[Who<sup>u</sup> left? (I don't know but) he<sub>u</sub> was furious.]]  
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- Wh-words are indefinites.
- Their inquisitiveness is introduced by the Foc head that they agree with.
- This does not affect their binding possibilities.
- Their licensing of anaphora is captured.

(28) [[Who<sup>u</sup> left? Was he<sub>u</sub> furious?]]  
= !([u]; left{u}); ?u; ?!(furious{u})

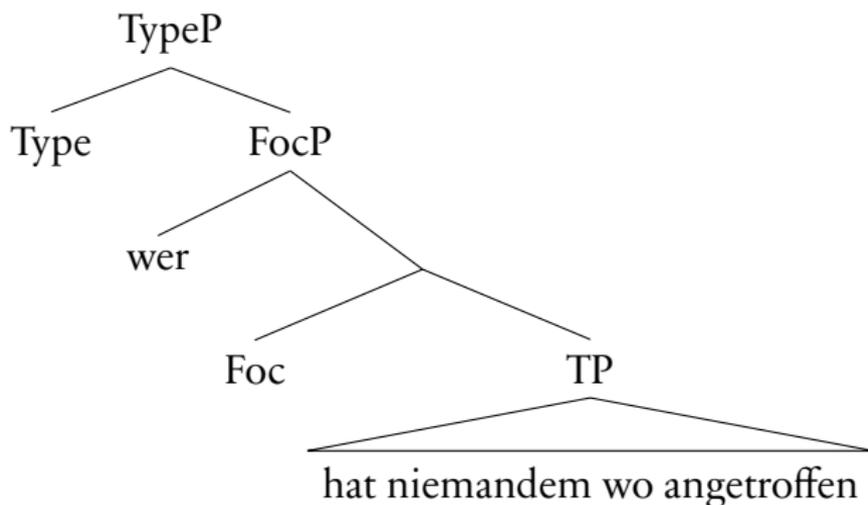
# Intervention effects

- Baseline: a grammatical example



- (29)  $\llbracket$  Who met Luise where? (German)  $\rrbracket$   
 $= !([u_1]; [u_2] \wedge meet(u_1, L, u_2)); ?u_1; ?u_1u_2$

## Intervention effects



(30)  $\llbracket$  Who met nobody where? (German)  $\rrbracket$   
 $= !([u_1]; \neg([u_3] \wedge [u_2] \wedge M(u_1, u_3, u_2))) ; ?u_1 ; ?u_1 u_2$

- Negation blocks discourse referents from projecting

## Intervention effects, II

- Mayr (2014): plural quantifiers intervene only when interpreted distributively

- (31) a. Wo haben sich mehr als drei Maler wann eine  
where have self more than three painters when a  
Pizza geteilt?  
pizza shared  
'Where did more than three painters share a pizza  
when?'
- b. \*Wo haben sich mehr als drei Maler wann eine  
where have self more than three painters when a  
Arbeitschuh angezogen?  
dungaree put.on

- Mayr (2014): plural quantifiers intervene only when interpreted distributively
- This contrast is predicted in our account because only distributive readings of quantifiers block the projection of dynamic information (cf. Kamp and Reyle 1993):

- (32) a. More than three students shared a pizza. It was tasty.  
b. More than three students sipped a coffee. #It was cold.

## ?u – ever visible?

Q-particles are ?u

- words that can function both as indefinites and as wh-words (quexistentials)
- a separate morpheme (a Q(uestion)-particle) obligatorily accompanies quexistentials and its syntactic position can disambiguate the reading

- (33) a. Daa \*(sá) aawaxáa i éesh?  
what Q he.ate.it your father  
‘What did your father eat?’ (Tlingit)
- b. Tlél goodéi \*(sá) xwagoot.  
not where.to Q I.went  
‘I didn’t go anywhere.’ (Tlingit)

Cable (2010)

## Q-particle – overt realization of the identity operator

- Q-particles can be separated from quexistentials
- Whenever a Q-particle can appear in two structurally different positions and the two positions differentiate between indefinite and wh-interrogative interpretations, then the low position must be the indefinite interpretation and the high position must be the wh-question interpretation.

Correct for Japanese, Sinhala, Tlingit

Hagstrom (1998), Cable (2010), Uegaki (2019)

# Non-exhaustive readings

Non-exhaustive readings straightforwardly captured:

(34) Who has a bike that I could borrow for 15 minutes?

	$w_a$	$w_{a,b}$	$w_b$	$w_\emptyset$
$u/a$	<input type="checkbox"/>		<input type="checkbox"/>	<input type="checkbox"/>
$u/b$	<input type="checkbox"/>	<input type="checkbox"/>		<input type="checkbox"/>

## Donkey anaphora in conditional questions

(35) If a farmer owns a donkey, does he beat it?

$$A_{(kk)} \twoheadrightarrow B_{(kk)} := \lambda c_k \lambda s_i. \quad s \in c \wedge \forall t \subseteq s$$

(36)  $(t \text{ subsists in } A(c))$   
 $\rightarrow t \text{ subsists in } B(A(c))$

$$\llbracket (35) \rrbracket = ([u_1]; [u_2]; \text{farmer}\{u_1\}; \text{donkey}\{u_2\}; \text{own}\{u_1, u_2\}) \\ \twoheadrightarrow \text{?!beat}\{u_1, u_2\}$$

# Exhaustive readings

To derive exhaustive readings, a **max** operator can be added to  $\text{Inq}_D$

$$(37) \quad \mathbf{max}\{u\} := \lambda c_k \lambda s_i. s \in c \wedge \forall p \in s \forall p' \in \bigcup c (w_p = w_{p'} \rightarrow g_{p'}(u) \leq g_p(u))$$

(38) Who did not register yet?

(39)  $!([u]; \text{notregistered}\{u\}; \mathbf{max}\{u\}); ?u$

	$w_a$	$w_{a,b}$	$w_b$	$w_\emptyset$
$u/a$				
$u/b$				
$u/ab$				

## Exhaustive readings and donkey anaphora

Brasoveanu (2008) derives weak / strong readings of donkey anaphora by postulating an ambiguity of indefinites (max  $\Rightarrow$  strong reading)

- (40)     a.    If a farmer owns a donkey, he beats it.  
          b.    If Bill has a dime, he puts it in the parking meter.

While singular indefinites are compatible with both readings, plural indefinites force strong readings:

- (41)     If Bill has dimes, he puts them in the parking meter.

The same effect is observed with exhaustive/non-exhaustive questions:

- (42)     Which students have a bike that I could borrow?

## Conclusion

- $\text{Inq}_D$  integrates insights from dynamic and inquisitive semantics in a single, compositional framework
- It is well-suited to capture both exhaustive and non-exhaustive question interpretations
- It can capture the anaphoric potential of wh-words
- It can be used to derive anaphora-related intervention effects
- It can connect the existence of exhaustive and non-exhaustive readings of questions to the existence of strong and weak readings of donkey anaphora.

THANK YOU

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## Appendix – lifting $\text{Inq}_D$ working with sets of assignments

# Dependencies

- Dependencies between assignments within one possibility:
- express quantificational dependency

(43) Each boy saw a movie. They liked it.

(44) Which boy saw which movie? Did they like it?

# Dependencies

- Dependencies between assignments within one possibility:
- express quantificational dependency
  - (43) Each boy saw a movie. They liked it.
  - (44) Which boy saw which movie? Did they like it?
- Dependencies between world and assignments:
- express uncertainty with respect to the value of the dref
  - (45) Someone dances.

## Extension and subsistence

A possibility  $\langle w', G' \rangle$  is an extension of another possibility  $\langle w, G \rangle$ ,  $\langle w', G' \rangle \geq \langle w, G \rangle$ , iff  $w' = w$  and  $G' \geq G$ .

(46) Dref matrix extension:

$$G' \geq G := \begin{cases} \forall g \in G. \exists g' \in G'. g \subseteq g' \wedge \\ \forall g' \in G'. \exists g \in G. g \subseteq g' \end{cases}$$

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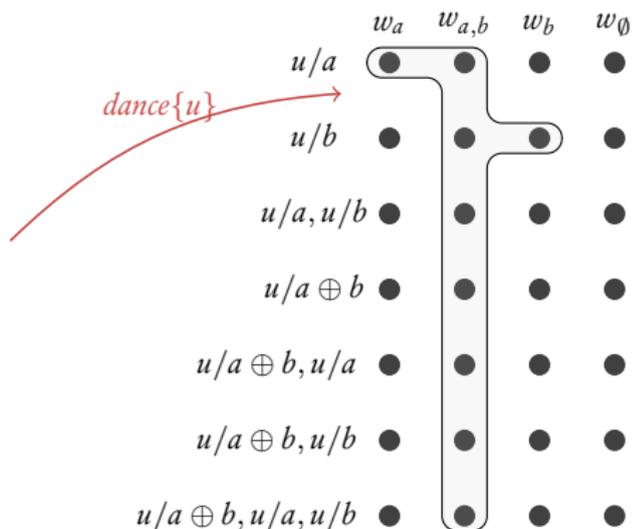
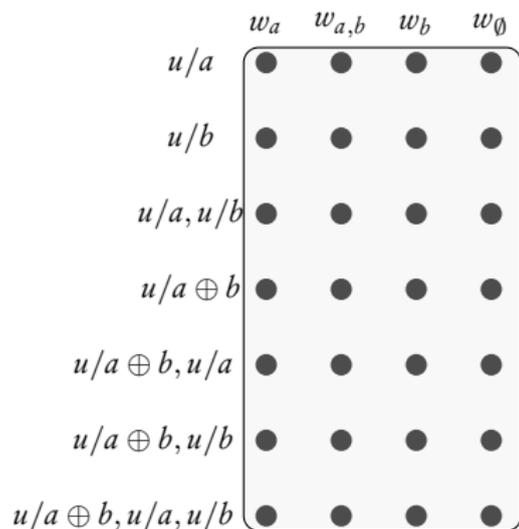
(47) State extension:  $s' \geq s := \forall p' \in s'. \exists p \in s. p' \geq p$

A state  $s$  **subsists** in a state  $s'$  iff:

- $s' \geq s$ , and
- $\forall p \in s. \exists p' \in s'. p' \geq p$

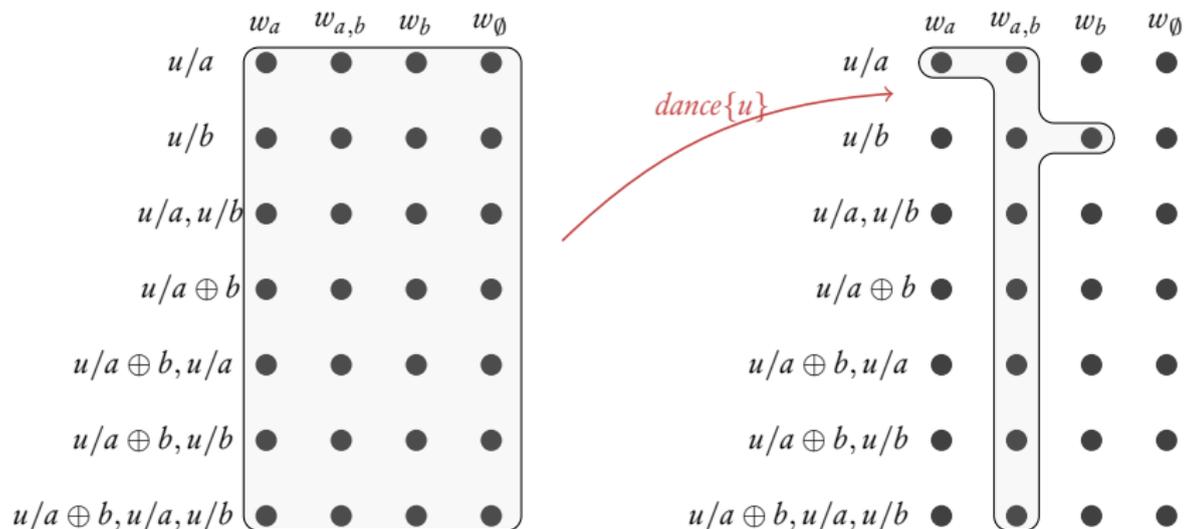
# Semantics: relations and conjunction

$$(48) \quad R\{u_1, \dots, u_n\} := \lambda c_k \lambda s_i. \quad s \in c \wedge \forall p \in s (\forall g \in G_p (R(w_p)(g(u_1)) \dots (g(u_n))))$$



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$$(49) \quad A_{(kk)}; B_{(kk)} := \lambda c_k. B(A(c))$$

## Semantics: dref introduction

- If  $C$  is a context and  $u$  not in assignments in  $C$ , then introducing the dref  $u$  creates the largest context  $C'$  such that every  $s' \in C'$  has some  $s \in C$  s.t.  $s$  subsists in  $s'$  and  $u \in \mathbf{assigns}(s')$ .

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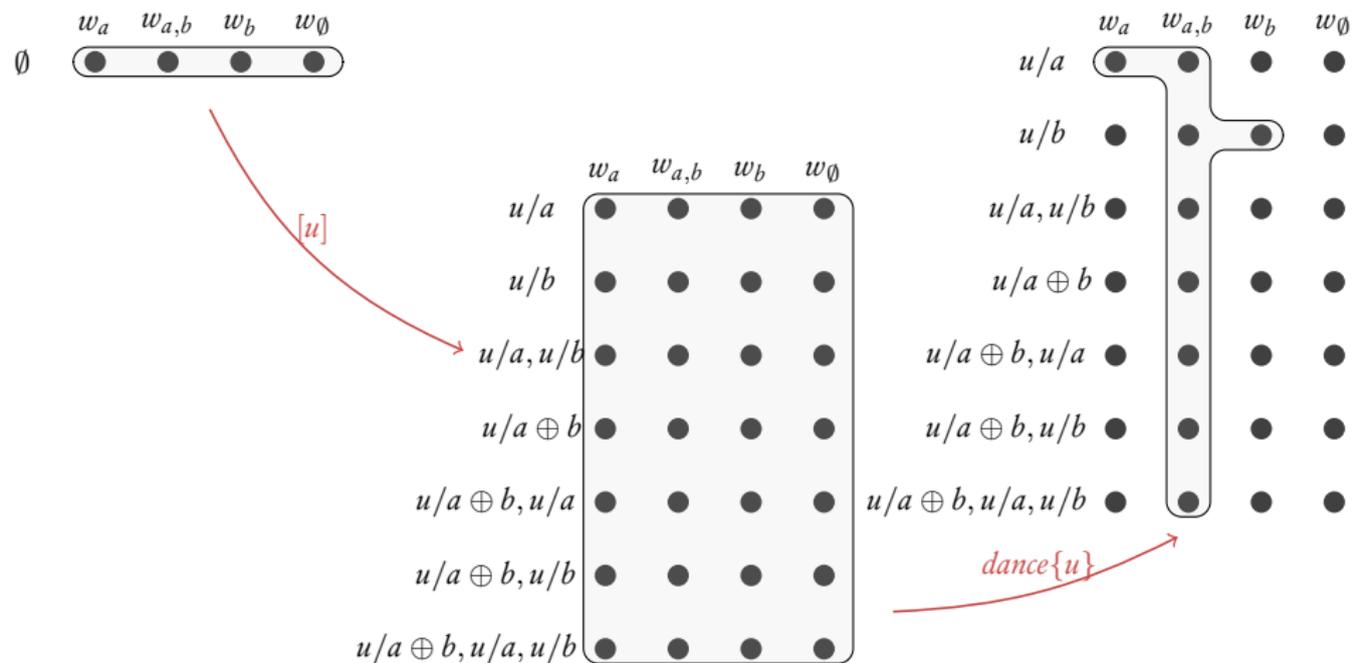
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- $\llbracket p[u]p' \rrbracket^{F,I,\theta} = 1$  iff  $\llbracket w_p = w_{p'} \rrbracket^{F,I,\theta} = 1$  and  $\llbracket G_p[u]G_{p'} \rrbracket^{F,I,\theta} = 1$

## Semantics: dref introduction

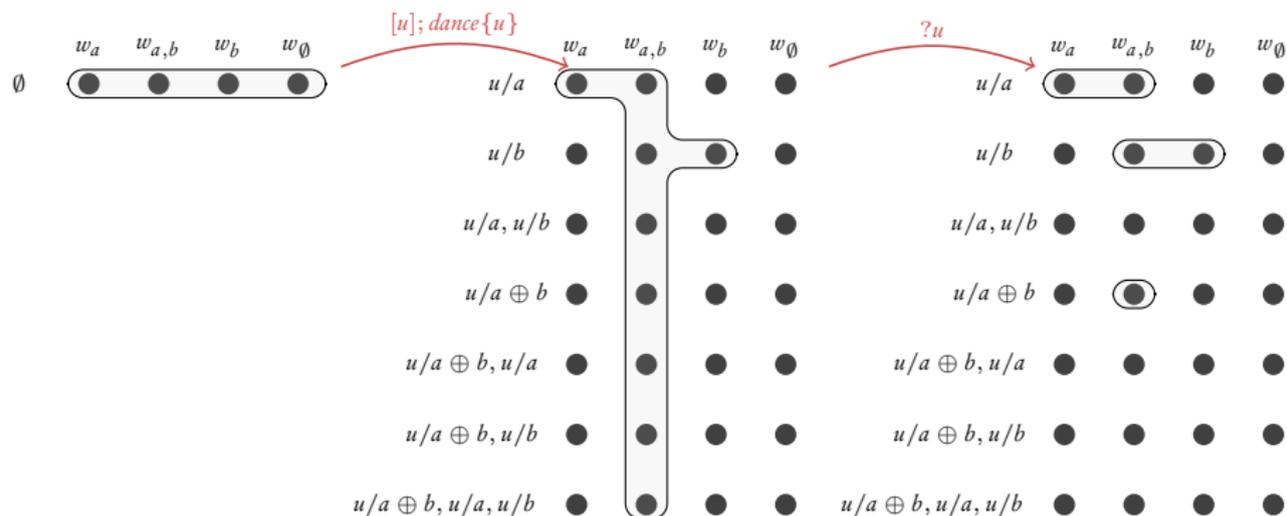
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- $[u] := \lambda c_k \lambda s_i. \left\{ \begin{array}{l} \exists s' \in c. \\ \forall p \in s. \exists p' \in s'. (p'[u]p) \wedge \\ \forall p' \in s'. \exists p \in s. (p'[u]p) \end{array} \right\}$

# Semantics: dref introduction



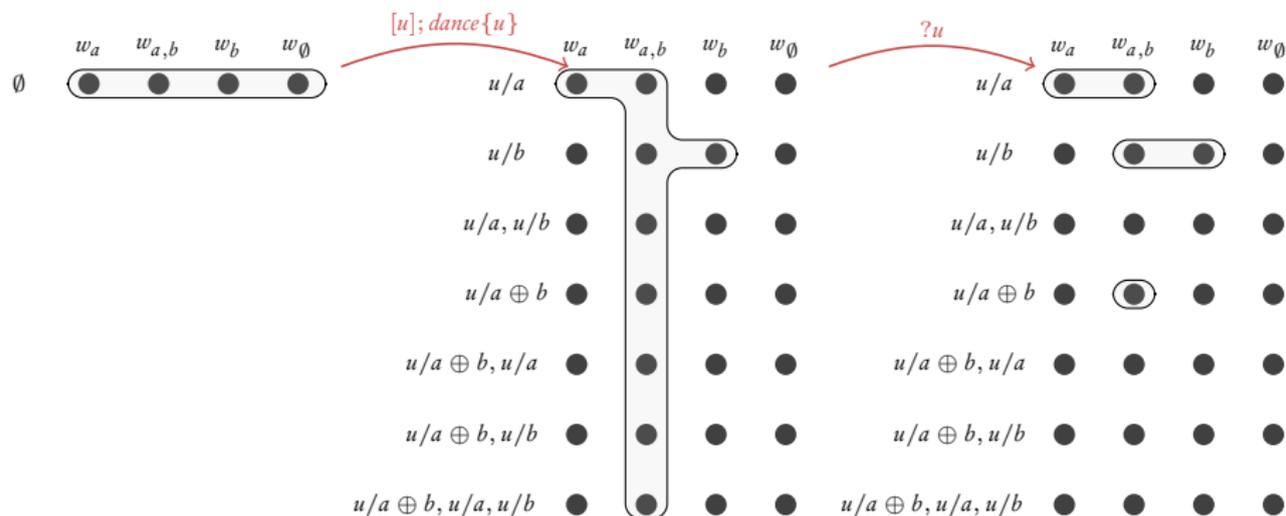
# Semantics: identity of drefs

$$(50) \quad ?u := \lambda c_k \lambda s_I. s \in c \wedge \exists x_e. \forall p \in s. \forall g \in G_p. x = g(u)$$



# Semantics: identity of drefs

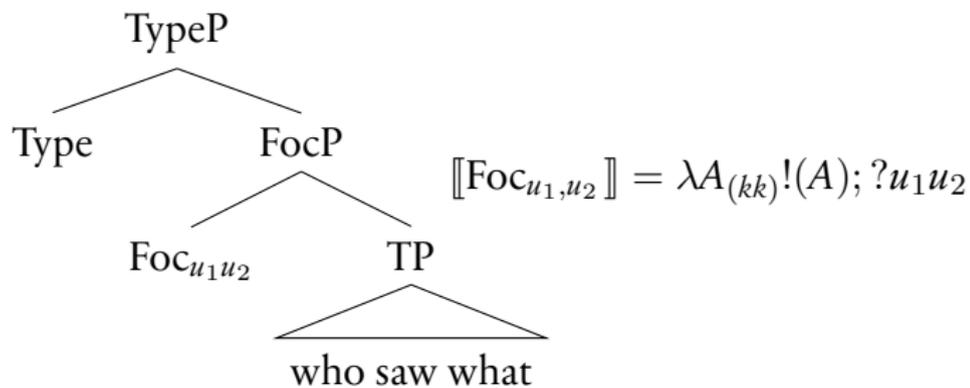
$$(50) \quad ?u := \lambda c_k \lambda s_I. s \in c \wedge \exists x_e. \forall p \in s. \forall g \in G_p. x = g(u)$$



(51)

$$?u_1, \dots, u_{n-1}, u_n := \lambda c_k \lambda s_I. s \in c \wedge \exists f(e^n, e). \forall p \in s. \forall g \in G_p. f(g(u_1), \dots, g(u_{n-1})) = g(u_n)$$

## Translations of wh examples



(52)  $\llbracket \text{Who saw what} \rrbracket = !([u_1]; [u_2]; \text{saw}\{u_1, u_2\}); ?u_1 u_2$

# Multiple wh-questions

(53)  $\llbracket \text{Who saw what} \rrbracket = !([u_1]; [u_2]; \text{saw}\{u_1, u_2\}); ?u_1u_2$

- The identification operator derives that pair-list reading is functional

(Higginbotham and May, 1981; Dayal, 1996)

- (54) Which student talked to which professor?
- Alice and Bill both talked to Prof. Carl.
  - #Alice talked to Prof. Carl and Prof. Dan.